## Indistinguishable Encryptions in the Presence of an Eavesdropper Class Exercise—2/22/18

Assume *G* is a PRG with input length *n* and output length n + 1. Do the following encryption schemes  $\Pi$  have indistinguishable encryptions in the presence of an eavesdropper? If yes, formally prove that if *G* is a PRG then the scheme is secure. If not, present a ppt adversary *A* and show that  $\Pr\left[\operatorname{PrivK^{eav}}_{A,\Pi}(n) = 1\right] \ge 1/2 + \rho(n)$  for some non-negligible  $\rho()$ .

1.  $\Pi$  is defined as follows: *Gen* outputs a random key k of length n. To encrypt a message  $m = m_1 || m_2$ , where  $m_1, m_2$  each have length n + 1, output  $c := (c_1 || c_2) := G(k) \oplus m_1 || G(k) \oplus m_2$ . To decrypt output  $m_1 || m_2 = G(k) \oplus c_1 || G(k) \oplus c_2$ .

2.  $\Pi$  is defined as follows: *Gen* outputs a random key k of length n. To encrypt a message m, where m has length n + 1, output  $c := G(k) \oplus m || 0^n$ . To decrypt, output the first n bits of  $c \oplus (G(k) || 0^n)$ .