ENEE 459E/CMSC 498R: Introduction to Cryptology
RSA Cryptanalysis
5/2/17

1. Partially Known Message.

Coppersmith's Theorem: Let $p(x)$ be a polynomial of degree e. Then in time $poly(\log(N), e)$ one can find all m such that $p(m) = 0 \ mod \ N$ and $m \leq N^{1/e}$.

Assume message is $m = m_1||m_2$, where m_1 is known, but m_2 (which consists of k bits) is not known. Using Coppersmith's Theorem, show how to recover m given the ciphertext c, assuming k is not too large.

Hint: Note that m can be expressed as $m \ := 2^k m_1 + m_2$.

2. Related Messages.

Euclidean Algorithm for Polynomials: Let $f(x)$ and $g(x)$ be two polynomials over $Z^*{}_N$. Then a slightly modified version of the Euclidean GCD Algorithm can be used to determine the greatest common divisor of f, g as polynomials over $Z^*{}_N$.

Assume the sender encrypts both m and $m + \delta$, for known δ, unknown m giving two ciphertexts c_1 and c_2. Use the Euclidean algorithm for polynomials to show how to recover m given knowledge of δ and given the two ciphertexts c_1, c_2.

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3. Sending the same message to multiple receivers:

The following is a slightly extended version of Chinese Remainder Theorem than the one we saw in class for the case where there are 3 moduli.

Chinese Remainder Theorem. Let N_1, N_2, N_3 be pairwise relatively prime. Then for every c_1, c_2, c_3, there exists a unique non-negative integer $\hat{c}$ such that:

$$\hat{c} = c_1 mod\ N_1$$
$$\hat{c} = c_2 mod\ N_2$$
$$\hat{c} = c_3 mod\ N_3.$$

Assume there are three receivers with public keys:
$pk_1 = \langle N_1, 3\rangle, pk_2 = \langle N_2, 3\rangle, pk_3 = \langle N_3, 3\rangle$.
A sender sends the same encrypted message m to all three receivers so an eavesdropper sees:
$c_1 = m^3\ mod\ N_1,\ c_2 = m^3\ mod\ N_2,\ c_3 = m^3\ mod\ N_3$

Show how to use the Chinese Remainder Theorem to recover m.