# Introduction to Cryptology 

Lecture 20

## Announcements

- HW8 due Tuesday, 4/19
- TA OH Monday 4/18 4pm-6pm


## Agenda

- More Number Theory!


## Generalized Theorem

Theorem: Let $G$ be a finite group with $m=|G|$, the order of the group. Then for any element $g \in G, g^{m}=1$.

Corollary of Fermat's Little Theorem is a special case of the above when $G$ is the multiplicative group $Z^{*}{ }_{p}$ and $p$ is prime.

## Multiplicative Groups Mod N

- What about multiplicative groups modulo $N$, where $N$ is composite?
- Which numbers $\{1, \ldots, N-1\}$ have multiplicative inverses mod $N$ ?
$-a$ such that $\operatorname{gcd}(a, N)=1$ has multiplicative inverse by Extended Euclidean Algorithm.
$-a$ such that $\operatorname{gcd}(a, N)>1$ does not, since $\operatorname{gcd}(a, N)$ is the smallest positive integer that can be written in the form $X a+Y N$ for integer $X, Y$.
- Define $Z^{*}{ }_{N}:=\{a \in\{1, \ldots, N-1\} \mid \operatorname{gcd}(a, N)=1\}$.
- $Z^{*}{ }_{N}$ is an abelian, multiplicative group.
- Why does closure hold?


## Order of Multiplicative Groups Mod N

- What is the order of $Z^{*}{ }_{N}$ ?
- This has a name. The order of $Z^{*}{ }_{N}$ is the quantity $\phi(N)$, where $\phi$ is known as the Euler totient function or Euler phi function.
- Assume $N=p \cdot q$, where $p, q$ are distinct primes.
$-\phi(N)=N-p-q+1=p \cdot q-p-1+1=$ $(p-1)(q-1)$.
- Why?


## Order of Multiplicative Groups Mod N

General Formula:
Theorem: Let $N=\prod_{i} p_{i}{ }^{e_{i}}$ where the $\left\{p_{i}\right\}$ are distinct primes and $e_{i} \geq 1$. Then

$$
\phi(N)=\prod_{i} p_{i}^{e_{i}-1}\left(p_{i}-1\right)
$$

## Another Special Case of Generalized Theorem

Corollary of generalized theorem:
For $a$ such that $\operatorname{gcd}(a, N)=1$ :

$$
a^{\phi(N)} \equiv 1 \bmod N
$$

## Another Useful Theorem

Theorem: Let $G$ be a finite group with $m=|G|>$ 1. Then for any $g \in G$ and any integer $x$, we have

$$
g^{x}=g^{x \bmod m}
$$

Proof: We write $x=a \cdot m+b$, where $a$ is an integer and $b \equiv x \bmod m$.

- $g^{x}=g^{a \cdot m+b}=\left(g^{m}\right)^{a} \cdot g^{b}$
- By "generalized theorem" we have that

$$
\left(g^{m}\right)^{a} \cdot g^{b}=1^{a} \cdot g^{b}=g^{b}=g^{x \bmod m}
$$

## An Example:

Compute $3^{25} \bmod 35$ by hand.

$$
\begin{gathered}
\phi(35)=\phi(5 \cdot 7)=(5-1)(7-1)=24 \\
3^{25} \equiv 3^{25} \bmod 24 \bmod 35 \equiv 3^{1} \bmod 35 \\
\equiv 3 \bmod 35 .
\end{gathered}
$$

## Background for RSA

Recall that we saw last time that

$$
a^{m} \equiv a^{m \bmod \phi(N)} \bmod N
$$

For $e \in Z^{*}{ }_{N}$, let $f_{e}: Z_{N}^{*} \rightarrow Z_{N}^{*}$ be defined as $f_{e}(x):=x^{e} \bmod N$.
Theorem: $f_{e}(x)$ is a permutation.
Proof: To prove the theorem, we show that $f_{e}(x)$ is invertible.
Let $d$ be the multiplicative inverse of $e \bmod \phi(N)$.
Then for $y \in Z_{N}^{*}, f_{d}(y):=y^{d} \bmod N$ is the inverse of $f_{e}$.
To see this, we show that $f_{d}\left(f_{e}(x)\right)=x$.
$f_{d}\left(f_{e}(x)\right)=\left(x^{e}\right)^{d} \bmod N=x^{e \cdot d} \bmod N=x^{e \cdot d \bmod \phi(N)} \bmod N=x^{1} \bmod N=$ $x \bmod N$.

Note: Given $d$, it is easy to compute the inverse of $f_{e}$
However, we saw in the homework that given only $e, N$, it is hard to find $d$, since finding $d$ implies that we can factor $N=p \cdot q$.
This will be important for cryptographic applications.

## Chinese Remainder Theorem

# Going from $(a, b) \in Z_{p} \times Z_{q}$ <br> $$
\text { to } x \in Z_{N}
$$ 

Find the unique $x \bmod N$ such that

$$
\begin{aligned}
& x \equiv a \bmod p \\
& x \equiv b \bmod q
\end{aligned}
$$

Recall since $\operatorname{gcd}(p, q)=1$ we can write

$$
X p+Y q=1
$$

Note that

$$
\begin{aligned}
& X p \equiv 0 \bmod p \\
& X p \equiv 1 \bmod q
\end{aligned}
$$

Whereas

$$
\begin{aligned}
& Y q \equiv 1 \bmod p \\
& Y q \equiv 0 \bmod p
\end{aligned}
$$

Going from $(a, b) \in Z_{p} \times Z_{q}$

$$
\text { to } x \in Z_{N}
$$

Find the unique $x \bmod N$ such that

$$
\begin{aligned}
& x \equiv a \bmod p \\
& x \equiv b \bmod q
\end{aligned}
$$

Claim:

$$
\begin{aligned}
& b \cdot X p+a \cdot Y q \equiv a \bmod p \\
& b \cdot X p+a \cdot Y q \equiv b \bmod q
\end{aligned}
$$

Therefore, $x \equiv b \cdot X p+a \cdot Y q \bmod N$

