

# ENEE/CMSC/MATH 456

## SPN Class Exercise

1. Present an attack and analyze the complexity of your attack to recover the all sub-keys of a two-round SPN (with a final key-mixing step) with the following parameters (same as picture on the attached sheet and the one in the lecture notes):
  - Block size:  $\ell = 16$
  - Sub-key length:  $n = 16$ , the three sub-keys,  $k_1, k_2, k_3$  are uniform, independent 16-bit keys.
  - Number of S-boxes: 4, each with 4-bit input/output  
Same structure as in the picture on the next sheet.

Solution: We brainstormed several solutions. The final one was as follows:

- Obtain an input/output pair  $(x,y)$
- Guess all possible  $k_3$  ( $2^{16}$  of them)
- Work backward to obtain the intermediate value after the  $k_2$  mixing step.
- Guess the values of  $k_2$  corresponding only to the outputs of the first S-box ( $2^4$  of them)
- Work backwards to obtain the intermediate value after the  $k_1$  mixing step.
- XOR with the appropriate bits of the input to obtain a candidate (partial)  $k_1$  value.
- We now have a table with  $(2^{20})$  candidate (partial) key tuples. We will ask for additional input/output pairs. Note that we have to work backwards from the output to obtain a partial input and that the partial input is 4-bit length. So we expect to require 5 additional input/output pairs as  $(2^4)^5 = 2^{20}$ .
- In total, we have spent  $5 \cdot 2^{20}$  time. We must repeat this 4 times to obtain the rest of the key. So this would be  $20 \cdot 2^{20} = 5 \cdot 2^{22}$  time.

Note this is still better than brute force search ( $2^{48}$  time) and better than our first attempt which did not make use of partitioning the input w.r.t S-boxes, which was  $2^{32}$ -time.

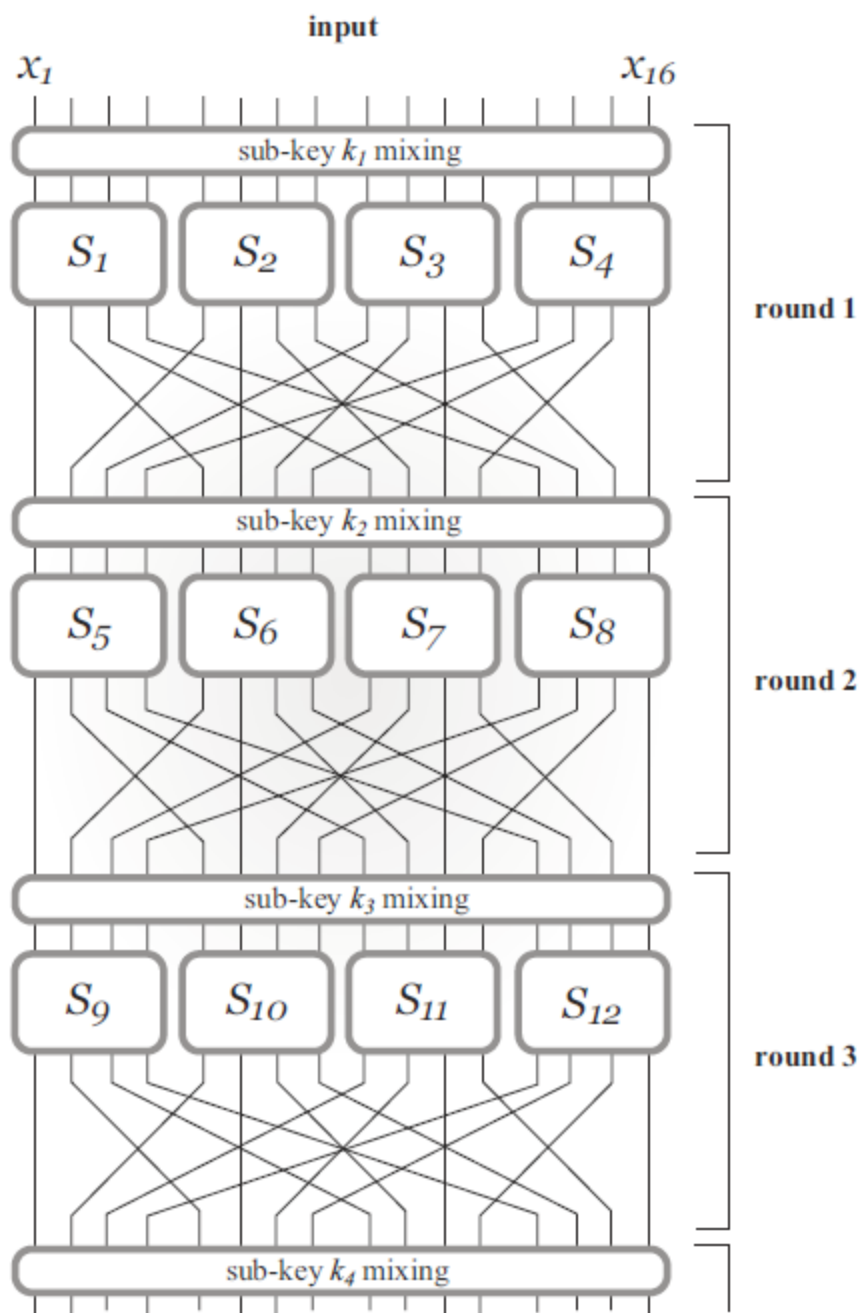


FIGURE 6.2: A substitution-permutation network.