All answers should be accompanied with proofs or sufficient explanation. Intermediate calculations should be shown.

In the problems you may use any representation of \mathbf{F}_{11} in the calculations. The *final answers* should be written using integers mod 11.

(a). (10pt). Prove that $\alpha = 2$ is a primitive element of $F = \mathbf{F}_{11}$.

(b). (10pts). Let C be a [10, 6, d] RS code over F. Write out a parity-check H matrix of C.

(c). (20pt) Reduce H to a systematic form $H' = [I_4|A]$.

Help: $\begin{bmatrix} 1 & 2 & 4 & 8 \\ 1 & 4 & 5 & 9 \\ 1 & 8 & 9 & 6 \\ 1 & 5 & 3 & 4 \end{bmatrix}^{-1} = \begin{bmatrix} 2 & 1 & 8 & 1 \\ 6 & 0 & 10 & 6 \\ 2 & 5 & 0 & 4 \\ 7 & 7 & 2 & 6 \end{bmatrix} \pmod{11}.$ (If you find no use of this equality, ignore it.)

(d). (30pt) Write out a generator matrix of C in a systematic form

(Use caution: Lecture 2 applies to binary codes only)

(e). (10pt) Using H', find the codeword c_0 that corresponds to the message symbols (1, 1, 1, 1, 1, 1).

(f). (10pt) What is the polynomial f such that $eval(f) = c_0$?

(g). (15pt) Is is true that $(c_0, c_1, ..., c_9) \in C$ implies that $(c_9, c_1, c_2, ..., c_8) \in C$?

(h). (20pt) Let $c \in C$ be a vector of weight 5. Prove that if $c' \in C$ is such that $\operatorname{supp}(c) = \operatorname{supp}(c')$ then c' = ac where $a \in F \setminus \{0\}$ is some constant.

(i). (15pt) Using problem 8, compute directly, with proof, the number of vectors of weight 5 in C (your answer should be a number, not an expression).

(j). (30pt) Let r = (3, 0, 0, 10, 5, 4, 0, 6, 10, 0) be a received vector. Perform the Peterson-Gorenstein Zierler algorithm to determine the number of errors and to decode the vector.

(k). (10pt) Explain how to find the polynomial f such that eval(f) equals the decoded codeword from problem 10.

(l). (10pt) If time left, invent a nice problem and solve it.